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performance software for SPECT  
imaging**

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# MEDITOMO\*

an high performance software package  
for 3D SPECT imaging

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## Abstract

We describe the parallel software library, named MEDITOMO, designed for analysis of MEDICAL images obtained by SPECT (*Single Photon Emission Computed Tomography*) TOMOgraphic systems. The library has been developed within a national research project carried out in collaboration with the Clinic Physiopathology Department of University of Florence, the Physics Department of University of Genoa and the Institute of High Performance Computing and Networking of CNR in Naples. MEDITOMO is the core library of the PSE (Problem Solving Environment) MEDIGRID, oriented to medical imaging applications, that will be employed in a grid infrastructure involving such departments.

## 1 The SPECT imaging model equation

The role of accurate and efficient investigation and diagnosis in the management of all disease is unquestionable. In last decades, central to the diagnostic process are imaging techniques, medical imaging not only provides for diagnosis but also serves to assist the treatment of malignant disease. Due to the huge amount of data, the exploitation of high performance environments both for the storage requirements and for their efficient elaboration, is mandatory. In particular, here we are concerned with the SPECT imaging, mathematical description of a SPECT analysis is based on the *Blurred Radon Transform* [12]:

$$\int_{-\infty}^{\infty} \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} u(\mathbf{x}, z') h(s - \mathbf{x} \cdot \theta, \mathbf{x} \cdot \theta^{\perp}, z - z') d\mathbf{x} dz' = g(s, \phi, z) \quad (1)$$

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where  $\mathbf{x} = \{x, y\}$ ,  $\theta = \{\cos \phi, \sin \phi\}$  and  $\theta^\perp = \{-\sin \phi, \cos \phi\}$ . If  $t = \mathbf{x} \cdot \theta^\perp$  is the distance between the source and the detector, coordinates  $\{s, t, z\}$  refer to the plane  $\{s, t\}$  which is rotated, with respect to the plane  $\{x, y\}$ , by the angle  $\theta$ .  $h(s, t, z)$  is the *Point Spread Function* (PSF) describing the *blurring* process, while  $g(s, \phi, z)$  represents the data acquired by the gamma camera at the angle  $\phi$  and  $u(x, y, z)$  refers to the radioisotope distribution.

Following [12] we assume:

$$h(s, t, z) = h_0(s, t)h_0(z, t) \quad ,$$

where:

$$h_0(a, b) = \frac{1}{\sigma(b)\sqrt{2\pi}}e^{-\frac{a^2}{2\sigma^2(b)}} \quad ,$$

and

$$\sigma(b) = \sigma_{min} - \frac{\sigma_{max} - \sigma_{min}}{2r'}(b - r') \quad , \quad (2)$$

$\sigma_{min}$ ,  $\sigma_{max}$  and  $r'$  depends on the acquisition system <sup>(1)</sup>. Equation (1) is said to be the *fully 3D* model. If the PSF is separable:

$$h(s, t, z) = E(s, t)A(z) \quad , \quad (3)$$

then:

$$\int_{-\infty}^{\infty} A(z - z')dz' \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} u(\mathbf{x}, z')h(s - \mathbf{x} \cdot \theta, \mathbf{x} \cdot \theta^\perp)dx = g(s, \phi, z) \quad . \quad (4)$$

Equation (4) is said to be the *2D+1* model [12].

Both the (1) and (4) can be written as:

$$g = Ku, \quad (5)$$

where  $K$  is the integral operator.

The mathematics of computed tomography involves the reconstruction of the unknown function  $u$  (usually a 3D surface) from its planar projections ( $g$ ) acquired at different values of  $\phi$ . More precisely, we deal with the inverse problem:

$$u = K^{-1}g \quad , \quad (6)$$

which is a well known inverse and ill posed problem [10]. Reconstruction methods available in literature can be largely classified into two groups, namely those based on *Filtered Back Projection* (FBP) and iterative methods. For a long time there was much debate as to the relative superiority of one or another and in particular whether one of the two was in some sense superior. Today, this debate has largely subsided with the inevitable conclusion that each method has its advantages, it being important to tailor the reconstruction algorithms to the

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<sup>1</sup>Some details about values of  $\sigma_{min}$ ,  $\sigma_{max}$  and  $r$  can be found in Section ??.

scanner design and to the physics of the imaging modality.

Both FBP and iterative methods can be related to regularization approaches. Basic idea of regularization is to find the solution  $u$  of (6) as:

$$u^* = \arg \min_u F(u) = \arg \min_u \{ \| Ku - g \|_2 \} \quad , \quad (7)$$

where  $\| \cdot \|_2$  is the norm defined on  $L^2(\Omega)$ ,  $\Omega \subset \mathbb{R}^3$  is an open subset of  $\mathbb{R}^3$  and is the domain of the  $u$  and  $g$  functions.

If the noise affecting data is assumed to have a poisson distribution, solution is find as:

$$u^* = \arg \max_u p(g|u) \quad , \quad (8)$$

where  $p(g|u)$  is the conditional probability<sup>2</sup> of  $g$  given  $u$  [26].

Both in (7) and in (8) we introduce as regularization functional the Total Variation seminorm, defined as:

$$TV(u) = \int |\nabla u| \, dx dy dz = \int \int \int \sqrt{u_x^2 + u_y^2 + u_z^2} \, dx dy dz \quad , \quad (9)$$

where  $u_x, u_y, u_z$  are partial derivatives of  $u(x, y, z)$  with respect to  $x, y$  and  $z$ . Then:

$$u^* = \arg \min_u F(u) = \arg \min_u \{ \| Ku - g \|_2^2 + \alpha TV(u) \} \quad (10)$$

and:

$$u^* = \arg \max_u \{ \log p(g|u) + TV(u) \} \quad . \quad (11)$$

In this paper we describe the design and development of the software library, named MEDITOMO, collecting parallel algorithms implementing the TV-based regularization for solving the SPECT imaging problem in high performance computing environment. The paper is organized as follows, we start from the description of the discretization of the problem in Section 2 and in Section 3 the basic algorithms underlying the parallel software are described: the Fixed Point, the Conjugate Gradient and the Expectation Maximization; in Section 4 the analysis of the computational complexity of the algorithms is carried out and in Section 5 how the parallelism has been introduced is discussed; in Section 6 main features of MEDITOMO are shown and, finally, in Section ?? some numerical results are shown.

## 2 Discretization

### 2.1 Discretization of $K$

Let:

$$s \in [-r, r] \quad , \quad z \in [-R, R] \quad , \quad \phi \in [0, 2\pi] \quad , \quad (12)$$

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<sup>2</sup>The maximum of  $p(g|u)$  can equivalently be found as the maximum of  $\log p(g|u)$

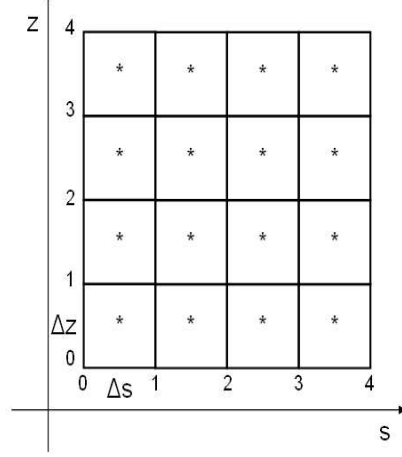


Figure 1: A  $4 \times 4$  grid for the discretization of  $g(s, \phi_m, z)$ . (\*) means the center of a cell.

and <sup>(3)</sup>:

$$\begin{aligned} \phi_m &= m\Delta\phi, & \Delta\phi &= \frac{2\pi}{M-1}, & m &= 0, \dots, M \\ s_k &= -r + k\Delta_s, & \Delta_s &= \frac{2r}{J-1}, & k &= 0, \dots, J-1 \\ z_l &= -R + l\Delta z, & \Delta z &= \frac{2R}{L-1}, & l &= 0, \dots, L-1 \end{aligned}$$

be the discretization of domain  $\Omega$ .

Let  $g^{kml}$  indicate the value of  $g(s, \phi, z)$  at  $(s_k, z_l, \phi_m)$ . Note that, at each  $m$ , and thus at each  $\phi_m$ , we have a 2D mesh made of  $J \times L$  boxes, as shown in figure 1. The centers  $C_{kml}$  of each square box have coordinates  $(s_k + \frac{\Delta s}{2}, \phi_m, z_l + \frac{\Delta z}{2})$ . Function  $u$  is discretized, by the same way, on a 3D mesh of size  $N \times N \times L'$ . We discretize the integrals in (1) using a midpoint quadrature formula which use the values of  $g$  at  $C_{kml}$ . Then, it is:

$$g^{kml} = \sum_{l'=0}^{L-1} \sum_{i=0}^{N-1} \sum_{j=0}^{N-1} k_{mkl}^{ijl'} \cdot u^{ijl'} \quad , \quad (13)$$

$$k = 0, \dots, J-1; \quad m = 0, \dots, M-1; \quad l = 0, \dots, L-1.$$

Equations (13) lead to the following linear system

$$[g] = [K]_3[u] \quad , \quad [g] \in \mathfrak{R}^{J \times M \times L} \quad , \quad [u] \in \mathfrak{R}^{N \times N \times L'} \quad (14)$$

where  $[u]$  is unknown.

<sup>3</sup>The values of  $J$ ,  $M$ ,  $L$  depend on the acquisition system. Typical values of  $M$  are 120, 60, and of  $J$  and  $L$  are  $2^l$ ,  $l = 7, 8$ .

**Definition 2.1** Matrix  $[K]_3 \in \mathfrak{R}^{((J \times M) \times L \times (N \times N) \times L')}$  is called the 3D projector. Elements of  $[K]_3$  are denoted by  $k_{mkl}^{ij}$ .  $[K]_3$  is a block matrix.

$[K]_3$ , is sparse with about 95% of zeros, then a compact storage technique was used.

Discretization of (4), leads to the following linear system:

$$g^{kml} = \sum_{i=0}^{N-1} \sum_{j=0}^{N-1} E_{km}^{ij} \sum_{l'=0}^{L'-1} A_{l'}^l \cdot u^{ijl'} \quad , \quad (15)$$

where

$$k = 0, \dots, J-1; \quad m = 0, \dots, M; \quad l = 0, \dots, L-1.$$

whose matrix formulation is:

$$[g] = [K]_3[u] \quad , \quad [K]_3 = [E]_2 \otimes [A]_1 \quad , \quad \text{where} \quad [E]_2 \in \mathfrak{R}^{((J \times M) \times (N \times N))} \quad (16)$$

$$\text{and} \quad [A]_1 \in \mathfrak{R}^{L \times L'} \quad , \quad (17)$$

where:

$$[E]_2 = \begin{pmatrix} E_{0,0}^{0,0} & \cdots & E_{0,0}^{0,N-1} & \cdots & E_{0,M-1}^{0,0} & \cdots & E_{0,M-1}^{0,N-1} \\ \vdots & \ddots & \vdots & \cdots & \vdots & \ddots & \vdots \\ E_{0,0}^{N-1,0} & \cdots & E_{0,0}^{N-1,N-1} & \cdots & E_{0,M-1}^{N-1,0} & \cdots & E_{0,M-1}^{N-1,N-1} \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ E_{J-1,0}^{0,0} & \cdots & E_{J-1,0}^{0,N-1} & \cdots & E_{J-1,M-1}^{0,0} & \cdots & E_{J-1,M-1}^{0,N-1} \\ \vdots & \ddots & \vdots & \cdots & \vdots & \ddots & \vdots \\ E_{J-1,0}^{N-1,0} & \cdots & E_{J-1,0}^{N-1,N-1} & \cdots & E_{J-1,M-1}^{N-1,0} & \cdots & E_{J-1,M-1}^{N-1,N-1} \end{pmatrix} \quad , \quad (18)$$

$[A]_1$  is a Toeplitz matrix:

$$[A]_1 = \begin{pmatrix} \alpha_0 & \cdots & \alpha_{L'-1} \\ \vdots & \ddots & \vdots \\ \alpha_{L-1} & \cdots & \alpha_0 \end{pmatrix} \quad , \quad (19)$$

where:

$$\alpha_l = (\Delta z)^2 \int_{-\frac{1}{2}}^{\frac{1}{2}} \int_{-\frac{1}{2}}^{\frac{1}{2}} A[\Delta z(\zeta - \zeta' + l)] d\zeta d\zeta' \quad , \quad (20)$$

**Definition 2.2**  $[E]_2 \in \mathfrak{R}^{(J \times M) \times (N \times N)}$  is called 2D projector.  $[E]_2$  is a block matrix.

Using the commutative property of the Kronecker product in (16) it follows:

$$[g] = [K]_3[u] = [E]_2 \otimes [A]_1[u] = [A]_1 \otimes [E]_2[u]$$

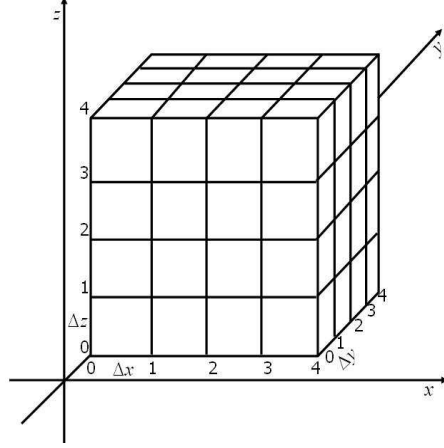


Figure 2: A  $4 \times 4$  grid used for the discretization of  $u(x, y, z)$

or:

$$\begin{aligned}
 [f] &= [E]_2[u] \quad , \\
 [g] &= [A]_1 \otimes [f] \quad . \\
 [u] &\in \mathfrak{R}^{(N \times N) \times L'} \quad , \quad [f] \in \mathfrak{R}^{(J \times M) \times L'} \quad , \quad [g] \in \mathfrak{R}^{(J \times M) \times L}
 \end{aligned} \tag{21}$$

## 2.2 Discretization of TV

The Euler Lagrange equations associated to the problem (10) are:

$$\begin{cases} K^*(Ku - z) + \alpha \cdot \mathcal{L}u = 0 & , \quad (x, y) \in \Omega \\ \frac{\partial u}{\partial n} = 0 & , \quad (x, y) \in \partial\Omega \end{cases} \tag{22}$$

where  $\mathcal{L}u = \text{div} \left( \frac{\nabla u}{|\nabla u|} \right)$

$\mathcal{L}$  is discretized, on the same mesh of  $u$  and  $g$  (see fig. 2), using a 7-point stencil (figura 3) [23], with step  $h = \Delta x = \Delta y = \Delta z$ . We approximate first and second partial derivatives of  $u(x, y, z)$  using the central finite differences:

$$\begin{aligned}
 \frac{\partial u}{\partial x} &\approx \frac{u_{i+1,j,l'} - u_{i-1,j,l'}}{2h} \quad , \\
 \frac{\partial u}{\partial y} &\approx \frac{u_{i,j+1,l'} - u_{i,j-1,l'}}{2h} \quad , \\
 \frac{\partial u}{\partial z} &\approx \frac{u_{i,j,l'+1} - u_{i,j,l'-1}}{2h} \quad ,
 \end{aligned}$$



$$\begin{aligned}\frac{\partial^2 u}{\partial x^2} &\approx \frac{u_{i+1,j,l'} + u_{i-1,j,l'} - 2u_{i,j,l'}}{h^2} , \\ \frac{\partial^2 u}{\partial y^2} &\approx \frac{u_{i,j+1,l'} + u_{i,j-1,l'} - 2u_{i,j,l'}}{h^2} , \\ \frac{\partial^2 u}{\partial z^2} &\approx \frac{u_{i,j,l'+1} + u_{i,j,l'-1} - 2u_{i,j,l'}}{h^2} .\end{aligned}$$

Discretization of  $\mathcal{L}$ , leads to :

$$L([u]) = B_h^T (D_h([u]))^{-1} B_h ,$$

where  $D_h$  is a block diagonal matrix:

$$D_h = \begin{bmatrix} D_x([u]) & 0 & 0 \\ 0 & D_y([u]) & 0 \\ 0 & 0 & D_z([u]) \end{bmatrix} ,$$

$B_h[u]$  is the matrix where each row contains just two not zero elements equal to 1.  $D_x([u])$ ,  $D_y([u])$ ,  $D_z([u])$  are diagonal matrices of size  $(N-1) \times N \times L'$ ,  $N \times (N-1) \times L'$ ,  $N \times N \times (L'-1)$  respectively.

Finally, discretization of (22) leads to the following non linear system:

$$([K^*]_3 [K]_3 - \alpha L([u]))[u] = [K^*]_3 [g] , \quad (23)$$

Equation (23) represents the discrete problem corresponding to the TV-based regularization problem describing the 3D SPECT imaging in case of additive gaussian noise. In the following we describe discretization of the same problem in case of a poisson noise.

### 2.3 Discretization of (8)

We perform the discretization of  $\Omega$  using the grid shown in fig. (2). Let  $u_s$  denote the values of  $u$  at the center of the  $s$ -th voxel, where  $s = 1, \dots, S$ . Consider the function  $\mathcal{J}$  defined as:

$$\mathcal{J} : (i, j, l') \mapsto s = (l' - 1) * N^2 + (j - 1) * N + i.$$

which maps  $(i, j, l')$  onto  $s$ ,  $s = 1, \dots, S$ . By the same way, we consider the function  $g(s, \phi, z)$  and suppose that the gamma camera acquires  $T$  values of  $g$ . Let denote by  $g_t, t = 1, \dots, T$  such values and:

$$u = \begin{pmatrix} u_1 \\ \vdots \\ u_S \end{pmatrix} , \quad (24)$$

$$g = \begin{pmatrix} g_1 \\ \vdots \\ g_T \end{pmatrix} . \quad (25)$$

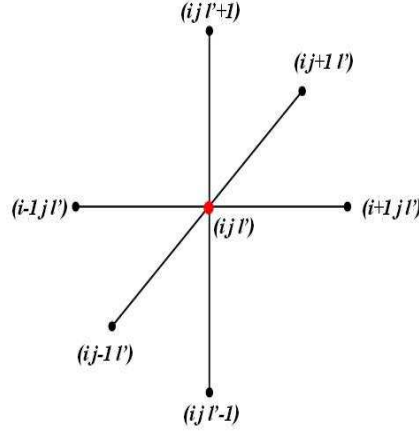


Figure 3: *The 7-point stencil used for discretization of  $\mathcal{L}$*

The difference between the discretization technique used in section 2.1 and this one is that in the first case function  $u$  is discretized at the points  $(i, j, l')$ , while now the values of  $u$  are computed at the centers of the voxels.

Let  $k_{ts}$  be the probability that the radio isotope radiation emitted from the generic voxel  $s = 1, \dots, S$  will be detected from the  $t$ -th ( $t = \dots, T$ ) detector, then:

$$g_t^* = \sum_{s=1}^S k_{ts} u_s = E(g|u) \quad (26)$$

is the average of the conditional statistic of  $g_t$ . The probability  $p(g_t|u)$  of  $g_t$  given  $u$  is:

$$p(g_t|u) = \text{Poisson} \left( \sum_{s=1}^S k_{ts} u_s \right) = e^{-\left(\sum_{s=1}^S k_{ts} u_s\right) g_t} \frac{\left(\sum_{s=1}^S k_{ts} u_s\right)^{g_t}}{g_t!} \quad (27)$$

Replacing (26) in the (27) gives:

$$p(g_t|u) = e^{-(g_t^*)^{g_t}} \frac{(g_t^*)^{g_t}}{g_t!} \quad (28)$$

then:

$$u^* = \arg \max_u \left\{ \sum_{t=1}^T \left( g_t \log \frac{g_t^*}{g_t!} - g_t^* \right) + \log p(u) \right\} \quad (29)$$

$u^*$  is solution of (??) if [35]:

$$u_s \frac{\partial l(u)}{\partial u_s} \Big|_{u_s=u_s^*} = 0, \quad u_s > 0 \quad (30)$$

$$\frac{\partial l(u)}{\partial u_s} \Big|_{u_s=u_s^*} \leq 0 \quad u_s = 0 \quad (31)$$

where

$$l(u) = \sum_{t=1}^T \left( g_t \log \frac{g_t^*}{g_t} - g_t^* \right) .$$

then:

$$0 = u_{s'} \frac{\partial l(u)}{\partial u_{s'}} = u_{s'} \left( - \sum_{t=1}^T k_{ts'} + \sum_{t=1}^T \frac{k_{ts'} g_t}{\sum_{s=1}^S k_{ts} u_s} \right), \quad (32)$$

or:

$$u_{s'} \sum_{t=1}^T k_{ts'} = u_{s'} \cdot \sum_{t=1}^T \frac{k_{ts'} g_t}{\sum_{s=1}^S k_{ts} u_s}. \quad (33)$$

We now introduce the matrix:

$$[K]_3 := (k_{ts})_{\substack{t=1, \dots, T \\ s=1, \dots, S}}$$

then (34) gives:

$$u([K]_3^T [1]) = u \left( [K]_3^T \frac{g}{([K]_3 u)} \right) \quad (34)$$

$$0 = -u([K]_3^T [1]) + u \left( [K]_3^T \frac{g}{([K]_3 u)} \right) \quad (35)$$

$$0 = -u \left\{ ([K]_3^T [1]) + \left( [K]_3^T \frac{g}{([K]_3 u)} \right) \right\} \quad (36)$$

where  $[1] \in \mathfrak{R}^t$  is a vector with all elements equal to 1, and where operations between vectors are intended component-wise.

### 3 The algorithms: Fixed Point and EM

We use two standard algorithms, the Fixed Point (FP) and the Expectation Minimization (EM), suitably modified for solving (23) and (36).

If  $\omega = \frac{\nabla[u]}{|\nabla[u]|}$ , (23) can be written as:

$$\begin{cases} \mathcal{F}(u, \omega) = [K^*]_3([K]_3[u] - [g]) + \alpha \nabla \cdot \omega = 0 \\ \omega = \frac{\nabla[u]_0}{\sqrt{(\nabla[u]_0)^2 + \beta^2}} \end{cases} \quad (37)$$

```

procedure PCG(K,K*,M,f,g);
* Computation of the residual at the zero step.
r(0) = g ;
* Preconditioning at the zero step.
z(0) = M-1 r(0);
* Search direction at the zero step.
p(0) = z(0);
γ0 = 0;
for k=1, maxit do
  * I. Computation of γk
  γk = ||r(k-1)||22;
  * II. Computation of βk
  βk = γk-1 / γk;
  * III. Update of the vector p(k)
  p(k) = r(k) + βkp(k-1)
  * IV. Update of the vector q(k)
  q(k) = K*M-1K p(k);
  αk = γk/||q(k)||22;
  * V. Update of the solution f(k)
  f(k) = f(k-1) + αkp(k);
  * VI. Update of the residual r(k)
  r(k) = r(k-1) - αkq(k);
endfor
end procedure PCG.

```

Figure 4: *PCG algorithm.*

Elimination of  $\omega$  from the first equation leads to:

$$[K^*]_3([K]_3[u] + [g]) + \alpha \nabla \cdot \left( \frac{\nabla[u]}{\sqrt{(\nabla[u])^2 + \beta^2}} \right) = 0 \quad . \quad (38)$$

If  $[u]_0 = [g]$  as in [13], we have:

$$[K^*]_3[K]_3[u]^{m+1} - [u]_0 + \alpha \nabla \cdot \left( \frac{\nabla[u]^{m+1}}{\sqrt{(\nabla[u]^m)^2 + \beta^2}} \right) = 0 \quad . \quad (39)$$

The iterative scheme (39) has a linear convergence [22].

The FP, shown in fig. 5, is made of two nested loops where the main operations are:

1. *I loop*: construction of the matrices in (39)
  - 1.a computation of  $[K]_3$ ;
  - 1.b computation of  $[K^*]_3$ ;
  - 1.c computation of  $L^m([u]^m)$ ;
2. *II loop*: solution of a linear system, by using the PCG :

The diagonal preconditioner is described in [29]. PCG is shown in fig. 4.

From (36) it follows the iterative scheme for computing the approximation  $u^*$  of (11) is:

$$u^{m+1} = \frac{[1]}{[K]_3^T [1]} u^m [K]_3^T \frac{g}{[K]_3 u^m} \quad m = 1, 2, \dots \quad (40)$$

which is known as the EM algorithm.

**Theorem 3.1** *The iterative scheme (40) converges to the solution  $u^*$  of the problem (??) [35].*

The explicit form of the EM algorithm is:

$$u_s^{m+1} = u_s^m \frac{1}{\sum_{t=1}^T k_{ts'}} \sum_{t=1}^T \frac{g_t k_{ts'}}{\sum_{s=1}^S k_{ts} u_s^m} \quad (41)$$

The EM (40) has a slow convergence rate [30], then the accelerated version, called EMOSn (*Expectation Maximization Ordered Subset*)[30] has been used. Let  $\Xi = \{\Omega_i : i = 1, \dots, n\}$  be a partition of the set  $\Omega = \{1, \dots, T\}$ , that is:

$$\begin{aligned} \Xi &= \{\Omega_i\}_{i=1, \dots, n}, \\ \Omega_i &\subseteq \Omega, \\ \bigcup_{i=1}^n \Omega_i &= \Omega, \\ \Omega_i \cap \Omega_j &= \emptyset, \quad \forall i \neq j \end{aligned} \quad (42)$$

```

begin procedure FP
{
  Construction of the matrix  $F$ ;
   $F = [K^*]_3[g]$ ;

  * Fixed Point Iteration

  for  $i = 1, \dots, \text{maxit}$  do
    Computation of the matrix  $L^m([u]^m)$ ;
    Solution of the system  $([K^*]_3[K]_3 + \alpha L^m([u]^m)) [u] = F$ 
    by the PCG method (nested cycle);
  endfor
}end procedure FP

```

Figure 5: *FP algorithm.*

The partition  $\Xi$  leads to a row-block formulation of  $[K]_3$ :

$$[K]_3 = \begin{bmatrix} [K_1]_3 \\ \vdots \\ [K_n]_3 \end{bmatrix}, \quad [K_i]_3 \in \mathbb{R}^{\frac{T}{n} \times S} \tag{43}$$

$[K_i]_3$  is a restriction of the operator  $[K]_3$  to  $\Omega_i$ . In the  $2D+1$  model, each block  $[K_i]_3$  can be factorized as (16):

$$[K_i]_3[u] = ([A]_1 \otimes [E]_2)_i[u] \quad i = 1, \dots, n \tag{44}$$

In fig. 6 we show the EMOS $n$  algorithm.

### 3.1 Introduction of Total Variation inside EM.

Introducing TV in (36) leads to the following scheme [34]:

$$u_s^{m+1} = u_s^m \frac{1}{\sum_{t=1}^T k_{ts} + \alpha (L(u^m) u^m)_s} \sum_{t=1}^T \frac{g_t k_{ts}}{\sum_{s=1}^S k_{ts} u_s^m} \quad m = 1, 2, \dots \tag{45}$$

TV-based EMOS $n$ , is shown in figure 7. EMOS $n$ -TV is made of two loops which perform the following operations:

1. *I loop (EM):* definition of matrices in (??):

**procedure** EMOS $n$

1.  $m = 0$
2. Let  $\hat{u}^m$  be a constant and positive value
3. repeat until to the  $\hat{u}^m$  convergences
  - (a)  $u^1 = \hat{u}^m$
  - (b)  $m = m + 1$
  - (c) for each OS level  $i = 1, \dots, n$

- i. compute the product  $\mu^i = [K_i]_3 u^i$

$$\mu_t^i = \sum_{s=1, \dots, S} k_{ts} u_s^i, \quad t \in T_i$$

- ii. compute  $u^{i+1} = \frac{[1]}{[K]_3^T [1]} u^i [K_i]_3^T \frac{g}{\mu^i}$

$$u_s^{i+1} = u_s^i \frac{1}{\sum_{t=1}^T k_{ts}} \sum_{t \in T_i} \frac{g_t k_{ts}}{\mu_t^i}, \quad s = 1, \dots, S$$

- (d)  $\hat{u}^m = u^{n+1}$

4.  $u = \hat{u}^m$

**end procedure** EMOS $n$

Figure 6: *The EMOS $n$  algorithm*

**procedure** EMOS $n$ -TV(input:  $K, g$  ; output:  $u$ )

1.  $m = 0$
2. Let  $\hat{u}^m$  be a constant and positive value
3. repeat until to the  $\hat{u}^m$  convergence

- (a)  $u^1 = \hat{u}^m$
- (b)  $m = m + 1$
- (c) computation of  $L^m(u^1)$
- (d) For each OS level  $i = 1, \dots, n$ 
  - i. compute the product  $\mu^i = [K_i]_3 u^i$

$$\mu_t^i = \sum_{s=1, \dots, S} k_{ts} u_s^i, \quad t \in T_i$$

- ii. computation of  $u^{i+1} = \frac{[1]}{[K]_3^T [1] + \alpha L^m(u^1) u^i} u^i [K_i]_3^T \frac{g}{\mu^i}$

$$u_s^{i+1} = u_s^i \frac{1}{\sum_{t=1}^T k_{ts} + (L^m(u^1) u^i)_s} \sum_{t \in T} \frac{g_t k_{ts}}{\mu_t^i}, \quad s = 1, \dots, S$$

- (e)  $\hat{u}^m = u^{n+1}$

4.  $u = \hat{u}^m$

**end procedure** EMOS $n$ -TV

Figure 7: TV-based EMOS $n$  algorithm.

- 1.a construction of  $[K]_3$ ;
- 1.b construction of  $[K^*]_3$ ;
- 1.c construction of  $L^m([u]^m)$ ;
2. II loop (OS):
  - 2.a computation of  $\mu^i = [K_i]_3 u^i$ ;
  - 2.b computation of  $L^m([u]^m) u^i$ .

## 4 Performance analysis

Main difficulty of the development of efficient algorithms is the capability to analyze and predict their performance. MEDITOMO library has been designed with the aim of providing the most efficient algorithms. To this aim, first step is



the performance analysis of the basic algorithms. We now estimate the flops of CG and EMOS<sub>n</sub> <sup>(4)</sup> considering the number of floating-point operations  $N_{cgXX}$  of one CG's step, and the number of floating-point operations  $N_{em2d+1}$  <sup>(5)</sup> of one EM's step. We will use the following notations:

<b>nsez</b>	$L$	number of sections
<b>kdim</b>	$J$	size of each projection
<b>nang</b>	$M$	number of angular views
<b>nsez</b>	$L'$	number of projections
<b>ndim</b>	$N$	size of each section
<b>nos</b>	$n$	number of <i>subset</i>
<b>kos</b>	$\frac{M}{n}$	number of angular views contained in each subset
<b>nmat</b>	-	size of working area

**nmat** is the number of points inside the square inscribed circle with diameter equals to **ndim**, then

$$\mathbf{nmat} \approx \pi \left( \frac{\mathbf{ndim}}{2} \right)^2 .$$

One step of CG requires:

1. 2 scalar products between vectors;
2. 3 *axpy* operations<sup>(6)</sup>;
3. two matrix-vector operations<sup>(7)</sup>.

In case of the fully 3D model, steps 1. and 2., are performed using vectors of length  $\mathbf{nmat} \times \mathbf{nsez}$ , than they require:

$$N_1 = 12 \times \mathbf{nmat} \times \mathbf{nsez} \quad \text{flops.} \quad (46)$$

Step 3. requires <sup>8</sup>:

$$N_2 = \mathbf{nang} \times \mathbf{nsez} \times (24 + 12 \times \mathbf{ndim} + 505 \times \mathbf{nmat} + 45 \times \mathbf{kdim}) \quad . \quad (49)$$

<sup>4</sup>Looking at the algorithms in fig. 5, it follows that the flops of one FP's step should be obtained summing the operations of one CG's iteration, the construction of  $L^m([u]^m)$  at each step  $m$ , the product  $L^m([u]^m)[u]^{m+1}$  at each step  $m$ . Numerical experiments show that the computational complexity of last two operations (matrix product and construction) is negligible. By the same way, the computational cost of one TV-based EMOS<sub>n</sub>'s step (see figure 7) is equivalent to one EMOS<sub>n</sub>'s step.

<sup>5</sup>Estimate is given only for the  $2D+1$  model.

<sup>6</sup>We denote by *axpy*, the operation

$$y \leftarrow y + \alpha x \quad .$$

<sup>7</sup>The cost of such operations is of  $\mathcal{O}((J \times M) \times L \times (N \times N) \times L')$ , but the sparsity reduces it to of about 95%.

<sup>8</sup>If we don't consider the sparsity of  $[K]_3$ , then the (49) should be:

$$\tilde{N}_2 = 2 \times \mathbf{nsez}^2 \times \mathbf{nang} \times \mathbf{kdim} \times \mathbf{ndim}^2 \quad . \quad (47)$$

and:

$$\begin{aligned}
N_{cg3d} &= 12 \times \mathbf{nmat} \times \mathbf{nsez} + \\
&+ \mathbf{nang} \times \mathbf{nsez} \times (24 + 12 \times \mathbf{ndim} + 505 \times \mathbf{nmat} + 45 \times \mathbf{kdim}) \quad .
\end{aligned} \tag{50}$$

In case of the 2D+1 model, steps at 1 and 2, are performed using vectors of length  $\mathbf{ndim}^2 \times \mathbf{nsez}$ , and they require:

$$N_1 = 12 \times \mathbf{ndim}^2 \times \mathbf{nsez} + 1 \quad \textit{flops}. \tag{51}$$

Step 3. requires [8]<sup>9</sup>:

$$N_2 = \mathbf{nsez} \times (\mathbf{nang} \times (106 \times \mathbf{ndim}^2 + 61 \times \mathbf{kdim})) \quad . \tag{53}$$

and:

$$\begin{aligned}
N_{cg2d+1} &= 12 \times \mathbf{ndim}^2 \times \mathbf{nsez} + 1 + \\
&+ \mathbf{nsez} \times (\mathbf{nang} \times (106 \times \mathbf{ndim}^2 + 61 \times \mathbf{kdim})) \quad .
\end{aligned} \tag{54}$$

#### 4.1 Floating Point operations in EMOS<sub>n</sub>.

At each iteration of EMOS<sub>n</sub>, the following operations are performed:

1. 2 scalar products between vectors;
2. 2 matrix-vector operations (<sup>9</sup>).

Using the scheme shown in figure 6, step 1. is performed **nos** times, with vectors of length  $\mathbf{ndim}^2 \times \mathbf{nsez}$ , then it requires:

$$N_1 = 2 \times \mathbf{ndim}^2 \times \mathbf{nsez} \quad \textit{flops}. \tag{55}$$

If we assume that  $\mathbf{nsez} = 25$ ,  $\mathbf{nang} = 120$ ,  $\mathbf{kdim} = \mathbf{ndim} = 128$  we have:

$$N_2 \sim 7 \times 10^9 \quad , \quad \widetilde{N}_2 = 335 \times 10^9 \iff \frac{\widetilde{N}_2}{N_2} = 2\% \tag{48}$$

(48) confirms that the cost of both matrix-vector operations, at each iteration, reduces of about 98%<sup>[7]</sup>.

<sup>9</sup>If we don't consider the sparsity of the matrix  $[K]_3$ , then the (53) should have the value of (47). If we suppose that  $\mathbf{nsez} = 25$ ,  $\mathbf{nang} = 120$ ,  $\mathbf{kdim} = \mathbf{ndim} = 128$ , we have:

$$N_2 \sim 4 \times 10^9 \quad , \quad \widetilde{N}_2 = 335 \times 10^9 \iff \frac{\widetilde{N}_2}{N_2} = 1,5\% \quad \textit{flops}. \tag{52}$$

52) confirms that cost of both matrix-vector operations, at one CG step, reduces of about 98.5%.

The matrix-vector product at 2. requires <sup>10</sup> the following computational cost [8]:

$$N_2 = 2 \times \text{nsez} \times (\text{nang} \times (48 \times \text{ndim}^2 + 31 \times \text{kdim})) \quad . \quad (57)$$

and:

$$\begin{aligned} N_{em2d+1} &= 2 \times \text{nos} \times \text{nsez} \times \text{ndim}^2 + \\ &+ 2 \times \text{nsez} \times (\text{nang} \times (48 \times \text{ndim}^2 + 31 \times \text{kdim})) \quad . \quad (58) \end{aligned}$$

If we compare (54) and (58), and suppose  $\text{nos} \approx 6$ , then  $N_{cg2d+1} \approx N_{em2d+1}$ .

## 5 Parallelism into 3D SPECT algorithms

Starting the the performance analysis, parallelism has been introduced to reduce the exexution time of the most expensive operation:

$$z = [K^*]_3 [K]_3 u \quad . \quad (59)$$

where (59) is the *projection/backprojection* operation.

### 5.1 Parallelization of the *fully 3D* model

The most suitable communication topology is a *1D periodic grid* of size  $1 \times N_{proc}$ , where  $N_{proc}$  is the number of processors. We use a *cyclic column block* distribution for  $[K]_3$  and a *cyclic row block* distribution for  $[K^*]_3$ <sup>(11)</sup> Each processor loads  $r = J \times M/N_{proc}$  rows and  $c = N \times N$  columns of  $[E_i]_2$  and loads  $r = N \times N$  rows and  $c = J \times M/N_{proc}$  columns of  $[E_i]_2^*$ . The operations in (59) are performed as follows:

1. *projection* :  $\mathbf{f}_p^i = \sum_{j=0}^{L'-1} [E_{(j+i-1, \text{mod } L')}]_2 \cdot \mathbf{u}^j, i = 0, \dots, L-1;$
2. *backprojection* :  $\mathbf{z}_p^i = \sum_{j=0}^{L'-1} [E_{(j+i-1, \text{mod } L')}^*]_2 \cdot \mathbf{f}_p^j, i = 0, \dots, L-1;$
3. *collective sum* :  $\mathbf{z}^i = \sum_{p=0}^{N_{proc}-1} \mathbf{z}_p^i, i = 0, \dots, L-1;$

In figure 8 we show the parallel algorithm.

<sup>10</sup>If we don't consider the sparsity of the matrix  $[K]_3$ , than the (57) should have the value of (47). If we suppose that  $\text{nsez} = 25$ ,  $\text{nang} = 120$ ,  $\text{kdim} = \text{ndim} = 128$ , we have:

$$N_2 \sim 2 \times 10^9 \quad , \quad \widetilde{N}_2 = 335 \times 10^9 \iff \frac{N_2}{\widetilde{N}_2} = 0,5\% \quad \text{flops.} \quad (56)$$

(56) confirms that computational cost of both matrix-vector operations, at one EMOS<sub>n</sub> iteration, reduces, of about 99.5%.

<sup>11</sup>For the matrix  $[K^*]_3$  can be used a block form as in (??).

```

procedure  $[K^*]_3[K]_3u$ 
Cycle on the number of processors.
  for  $p = 0, N_{proc} - 1$ 
    Compute the partial product:  $\mathbf{f}_p^i \in \mathfrak{R}^{J \cdot M / N_{proc}}$ 
    step 1. for  $i = 0, \dots, L$ 
       $\mathbf{f}_p^i = E_i(p \cdot (J \cdot M / n_{proc}) : (p + 1) \cdot (J \cdot M / N_{proc}) - 1; :) \cdot \mathbf{u}^i;$ 
    Compute the partial product:  $\mathbf{z}_p^i \in \mathfrak{R}^{N \cdot N}$ 
    step 2. for  $i = 0, \dots, L$ 
       $\mathbf{z}_p^i = E_i^*(:, p \cdot (J \cdot M / n_{proc}) : (p + 1) \cdot (J \cdot M / N_{proc}) - 1) \cdot \mathbf{f}_p^i;$ 
    endfor
Collective sum operation.
 $\mathbf{z}^i = \sum_{p=0}^{N_{proc}-1} \mathbf{z}_p^i$ 
end procedure

```

Figure 8: *Parallel algorithm for  $z = [K^*]_3[K]_3u$  in the fully 3D problem.*

## 5.2 Parallelization of the $2D+1$ model

The operation (59) has a block formulation also in the  $2D+1$  model. Indeed, from the (15) each component  $g^l$  of  $[g]$ , of size  $J \times M$ , can be computed performing  $L'$  products using  $[E]_2$  and summing the resulting vectors multiplied by the  $\alpha_{l'-l}$ , that is:

$$\begin{pmatrix} \mathbf{g}^0 \\ \mathbf{g}^1 \\ \vdots \\ \mathbf{g}^{L'-1} \end{pmatrix} = \begin{pmatrix} \alpha_0[E]_2 & \alpha_1[E]_2 & \dots & \alpha_{L'-1}[E]_2 \\ \alpha_{-1}[E]_2 & \alpha_0[E]_2 & \dots & \alpha_{L'-1}[E]_2 \\ \vdots & \vdots & \ddots & \vdots \\ \alpha_{1-L}[E]_2 & \alpha_{2-L}[E]_2 & \dots & \alpha_0[E]_2 \end{pmatrix} \begin{pmatrix} \mathbf{u}^0 \\ \mathbf{u}^1 \\ \vdots \\ \mathbf{u}^{L'-1} \end{pmatrix}. \quad (60)$$

If the partition (42) holds then the (60) is valid for each  $[K_i]_3$ , that is:

$$\begin{pmatrix} (\mathbf{g}^0)_i \\ (\mathbf{g}^1)_i \\ \vdots \\ (\mathbf{g}^{L'-1})_i \end{pmatrix} = \begin{pmatrix} \alpha_0([E]_2)_i & \alpha_1([E]_2)_i & \dots & \alpha_{L'-1}([E]_2)_i \\ \alpha_{-1}([E]_2)_i & \alpha_0([E]_2)_i & \dots & \alpha_{L'-2}([E]_2)_i \\ \vdots & \vdots & \ddots & \vdots \\ \alpha_{1-L'}([E]_2)_i & \alpha_{2-L'}([E]_2)_i & \dots & \alpha_0([E]_2)_i \end{pmatrix} \begin{pmatrix} \mathbf{u}^0 \\ \mathbf{u}^1 \\ \vdots \\ \mathbf{u}^{L'-1} \end{pmatrix} \quad i = 1, \dots, n. \quad (61)$$

We introduce two parallelization approaches for computing the product (61).

```

procedure  $[K_i]_3 \mathbf{u}$ 
  step 1. for  $j = p * L_{loc} + 1, \dots, (p + 1) * L_{loc}$ 
     $p = ([E]_2)_k \mathbf{f}_j$ 
    for  $i = 0, 1, 2, \dots, L - 1$ 
       $(\mathbf{g}^i)_k = (\mathbf{g}^i)_k + \alpha_{j-i} p$ 
    endfor
  endfor
  step 2. for  $i = 0, 1, 2, \dots, L - 1$ 
    Collective sum operation
     $\mathbf{g}^i = \sum_{p=0}^{N_{proc}} (\mathbf{g}^i)_p$ 
  end procedure

```

Figure 9: Parallel algorithm for  $\mathbf{g} = [K_i]_3 \mathbf{u}$  in the 2D+1 model (First Strategy).

### 5.2.1 First strategy

Each processor  $p$  loads  $L_{loc} = L/N_{proc}$  blocks of columns of the matrix  $[K_i]_3$ , then operation (61) requires:

1. *projection:*  $((\mathbf{g}^j)_i)_p = \sum_{k=p*L'_{loc}+1, \dots, (p+1)*L'_{loc}} \alpha_{j-k} ([E]_2)_i \mathbf{u}^k \quad j = 0, \dots, L - 1;$
2. *sum:*  $\mathbf{g}^j = \sum_{p=0}^{N_{proc}} ((\mathbf{g}^j)_i)_p \quad , \quad j = 0, \dots, L - 1.$

In figure 9 we show the algorithm. The parallel algorithm for computing  $\mathbf{u} = [K_i^*]_3 \mathbf{g}$  is designed analogously.

### 5.2.2 Second strategy

Processor  $p$  loads  $L'_{loc} = L'/N_{proc}$  vectors  $\mathbf{u}^j$ ,  $L_{loc} = L/N_{proc}$  vectors  $(\mathbf{g}^i)_k$  and  $L'_{loc}$  blocks of columns of  $[K_i]_3$ . For each pair  $(p, q)$  of processors we introduce the set:

$$E_p^q = \left\{ i \in I^q : A_i^j \neq 0 \quad , \forall j \in J^p \right\}, \quad (62)$$

where  $I^p$  and  $J^p$  denote the set of the indices of the vectors  $(\mathbf{g}^i)_k$  and  $\mathbf{u}^j$  respectively assigned to processor  $p$ . Then  $E_p^q$  contains the indices of  $(\mathbf{g}^i)_k$ , which is assigned to processor  $q$ . Vectors  $\mathbf{u}^j$ , given to processor  $p$  contribute to the the computation of  $(\mathbf{g}^i)_k$ .

Computation of (61), can be performed by the following two steps:

1. *projection:*

- 1.a *Computation of  $(\mathbf{g}^i)_k$  of  $I^p$ :*

$$(\mathbf{g}^i)_k = \sum_{j \in J^p} \alpha_{i-j} ([E]_2)_k \mathbf{u}^j, \quad \forall i \in I^p$$

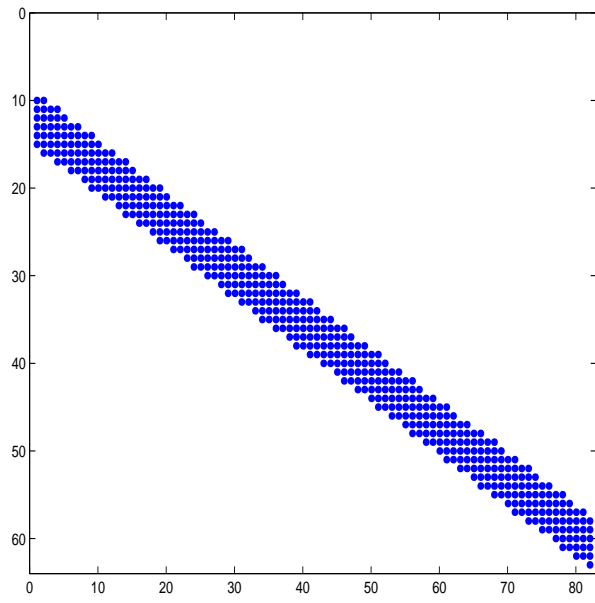


Figure 10: Representation of the sparsity pattern of the matrix  $[A]_1$ , with  $L = 68$  and  $L' = 82$ .

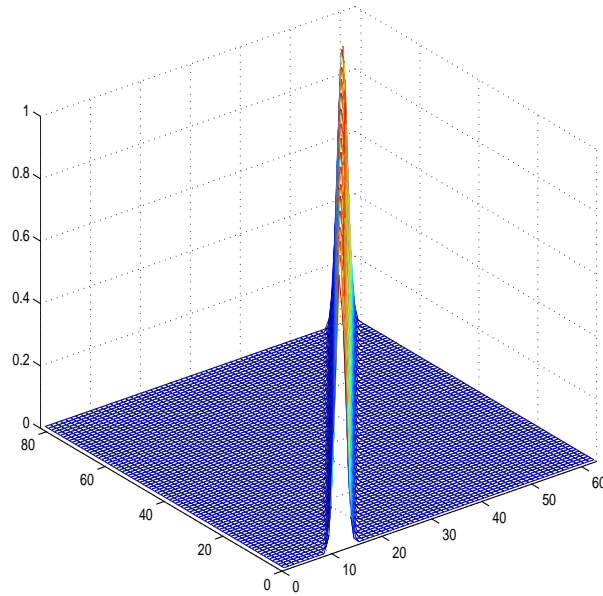


Figure 11: Elements of  $[A]_1$ , with  $L = 68$  and  $L' = 82$ .

1.b Contribution to  $(\mathbf{g}^i)_k$  by  $E_p^q$ :

$$((\mathbf{g}^i)_k)_p = \sum_{j \in J^p} \alpha_{i-j} ([E]_2)_k \mathbf{u}^j, \quad \forall q : E_p^q \neq \emptyset$$

2. sum between the neighborhood processors:

2.a send  $((\mathbf{g}^i)_k)_p$  to processor  $q$ ,  $\forall i \in E_p^q$ .

2.b receive  $((\mathbf{g}^i)_k)_q$  from processor  $q$ ,  $\forall q : E_q^p \neq \emptyset$

2.c local sum:

$$(\mathbf{g}^i)_k = (\mathbf{g}^i)_k + \sum_{q: E_q^p \neq \emptyset} ((\mathbf{g}^i)_k)_q, \quad \forall i \in E_q^p$$

In figure 12 we describe the parallel algorithm.

Parallel algorithm for computing  $\mathbf{u} = [K_i^*]_3 \mathbf{g}$  has been designed by the same way.

### 5.2.3 Performance comparison

We compare the strategies described above in terms of  $N_{calc}$ ,  $N_{com}$  and speed-up. Processors are logically organized in a 2D mesh of size  $1 \times N_{proc}$ . In step 1 of the first strategy, each processor  $p$  contributes to the computation of the  $L$  components of  $\mathbf{g}$  (or equivalently of  $(\mathbf{g}^i)_k$ ,  $i = 1, \dots, L$ ). This computation requires a block matrix-vector which requires  $\mathcal{O}((L \times L_{loc}) \times (J \times M) \times (N \times N))$  flops<sup>(12)</sup>, then the computational cost is:

$$N_{calc}^{(1.First)} = \mathcal{O}(N_{proc} \times (L \times L_{loc}) \times (J \times M) \times (N \times N)) \quad . \quad (63)$$

Step 2 requires a *collective sum* which requires<sup>(13)</sup>

$$N_{calc}^{(2.First)} = \mathcal{O}(N_{proc} \times \log_2(N_{proc}) \times (J \times M) \times L) \quad flops, \quad (64)$$

and uses  $N_{com}^{First}$  communications<sup>(14)</sup> where:

$$N_{com}^{First} = \mathcal{O}(N_{proc} \times \log_2(N_{proc}) \times (J \times M) \times L) \quad , \quad (65)$$

because the  $N_{proc}$  processors, in couples, send one to each other  $((\mathbf{g})_k)_p$  vectors of length  $L \times (J \times M)$ .

---

<sup>12</sup>Matrix-vector product uses local matrix and vector of size, respectively,  $L \times (J \times M) \times (N \times N) \times L_{loc}$  and  $(N \times N) \times L_{loc}$ .

<sup>13</sup>We assume that collective sum between vectors  $(\mathbf{g})_p$  (of length  $(J \times M) \times L$ ) is performed by  $N_{proc}$  processors using a tree communication topology where with depth  $\log_2 N_{proc}$ . At each level processors, organized in couples, communicate each other and update, using a sum operation, own vector  $(\mathbf{g})_p$ . Therefore, at each level,  $N_{proc}$  sum between vectors of length  $(J \times M) \times L$  are performed (MPI function `MPI_ALLREDUCE` used to perform the collective sum uses such topology)

<sup>14</sup>At each level  $N_{proc}$  communications are performed, therefore the total number of communications performed during the collective sum operation is  $N_{proc} \times \log_2 N_{proc}$ .

```

procedure  $([K]_3)_k f (\dots)$ 
  for each  $j \in J^p$ 
  begin for each
    Computation of the 2D projections
     $\mathbf{tu} = ([E]_2)_k \mathbf{u}_j$ 
    Computation of contribution of the weights
     $\alpha_{i-j}$  to projections
    for each  $i \in I^p$ 
    begin for each
       $(\mathbf{g}^i)_k = (\mathbf{g}^i)_k + \alpha_{i-j} \mathbf{tu}$ 
    end for each
    end for each
    send operation involving  $((\mathbf{g}^i)_k)_p$ 
    for each  $q \in \{1, \dots, N_{proc}\}$  such that  $E_p^q \neq \emptyset$ 
    begin for each
      for each  $j \in J^p$ 
      begin for each
         $\mathbf{tu} = ([E]_2)_k \mathbf{f}_j$ 
        for each  $i \in E_p^q$ 
        begin for each
           $((\mathbf{g}^i)_k)_p = ((\mathbf{g}^i)_k)_p + \alpha_{j-i} \mathbf{tu}$ 
          call  $\text{Send}(((\mathbf{g}^i)_k)_p, q)$ 
        end for each
      end for each
    end for each
  end for each
  Receive operations involving  $((\mathbf{g}^i)_k)_q$ 
  for each  $q \in \{1, \dots, N_{proc}\}$  such that  $E_q^p \neq \emptyset$ 
  begin for each
    for each  $i \in E_q^p$ 
    begin for each
      call  $\text{Receive}(((\mathbf{g}^i)_k)_q, q)$ 
      Local sum operation
       $(\mathbf{g}^i)_k = (\mathbf{g}^i)_k + ((\mathbf{g}^i)_k)_q$ 
    end for each
  end for each
end procedure

```

Figure 12: *Parallel algorithm for computing  $\mathbf{g} = [K_i]_3 \mathbf{u}$  of  $2D+1$  problem (Second Strategy).*



At first step of the second strategy, each processor  $p$  calculates  $|J^p|$  <sup>(15)</sup> components of  $\mathbf{g}$ , i.e. calculates  $(\mathbf{g}^i)_k$  where  $i \in J^p$ , and contributes to the computation of the other  $|E_p^q|$  components of  $\mathbf{g}$ , i.e. to the other projections  $(\mathbf{g}^i)_k$  where  $i \in E_p^q$ , by using a block matrix-vector operations, with a cost of  $\mathcal{O}((|J^p| + |E_p^q|) \times (J \times M) \times (N \times N))$  flops.  $|J^p| = L_{loc}$ , then the cost is:

$$N_{calc}^{(1.Second)} = \mathcal{O}(N_{proc} \times ((L_{loc} + |E_p^q|) \times L'_{loc}) \times (J \times M) \times (N \times N)) \quad \text{flops.} \quad (66)$$

Second step of the second strategy requires a sum between at most  $|E_p^q|$  neighborhood processors with a cost of:

$$N_{calc}^{(2.Second)} = \mathcal{O}(N_{proc} \times |E_p^q| \times (J \times M) \times |E_p^q|) \quad \text{flops,} \quad (67)$$

and a number of communications:

$$N_{com}^{(Second)} = \mathcal{O}(N_{proc} \times |E_p^q| \times (J \times M) \times |E_p^q|) \quad , \quad (68)$$

because, at each communication, processors  $p$  and  $q$  where  $|E_p^q| \neq 0$ , exchange sub-vectors of length  $|E_p^q| \times (J \times M)$  of  $(g)_k$ .

If  $L = L'$ , that is  $L_{loc} = L'_{loc}$ :

$L_{loc}$  is at most  $|E_p^q| = |J^q|$  and:

$$L \times L_{loc} = \frac{L^2}{N_{proc}} > 2 \frac{L^2}{N_{proc}^2} = (L_{loc} + |E_p^q|) \times L_{loc} \quad ,$$

then:

$$\boxed{N_{calc}^{(1.First)} > N_{calc}^{(1.Second)}} \quad (69)$$

If  $N_{proc}$  increases also  $\log_2 N_{proc}$  increases, and  $L_{loc}$  decreases, then:

$$\boxed{N_{calc}^{(2.First)} > N_{calc}^{(2.Second)}} \quad (70)$$

From  $N_{calc}^{(2.First)} = N_{com}^{(First)}$  and  $N_{calc}^{(2.Second)} = N_{com}^{(Second)}$ , and from (70) it follows:

$$\boxed{N_{com}^{(First)} > N_{com}^{(Second)}} \quad (71)$$

Let us now consider the speed-up  $S_{N_{proc}}$ :

$$S_{N_{proc}} = \frac{T_1}{T_{N_{proc}}} = \frac{N \times t_{calc}}{N_{calc} \times t_{calc} + T_{com} \times t_{com}} \quad , \quad (72)$$

---

<sup>15</sup> $|J|$  denotes the cardinality of the set  $J$ .

where  $T_1$  and  $T_{N_{proc}}$  denote the execution time respectively on one and  $N_{proc}$  processors,  $N$  and  $N_{calc}$  denotes the number of floating point operations performed respectively by one and by  $N_{proc}$  processors,  $t_{calc}$  is the execution time of one floating point operation, and  $t_{com}$  is the execution time of one communication between two processors.

The Product (61) requires:

$$N_{calc} = \mathcal{O}((L \times L) \times (J \times M) \times (N \times N)) \quad flops, \quad (73)$$

therefore:

$$T_1 = [(L \times L) \times (J \times M) \times (N \times N)] \times t_{calc} \quad . \quad (74)$$

If, for the sake of simplicity, we assume  $J = M = N$  then (74) can simplified as follow:

$$T_1 = [L^2 \times N^4] \times t_{calc} \quad . \quad (75)$$

and:

$$T_{N_{proc}}^{(*)} = [N_{calc}^{(1,*)} + N_{calc}^{(2,*)}] \times t_{calc} + [N_{com}^{(*)}] \times t_{com} \quad , \quad (76)$$

where the symbol “\*” means “*First*” or “*Second*”.

Substituting (63), (64) and (65) in (76) gives:

$$\begin{aligned} T_{N_{proc}}^{(First)} &= [N_{proc} \times (L \times L_{loc}) \times (J \times M) \times (N \times N) + \\ &+ N_{proc} \times \log_2(N_{proc}) \times (J \times M) \times L] \times t_{calc} + \\ &+ [N_{proc} \times \log_2(N_{proc}) \times (J \times M) \times L] \times t_{com} \quad . \end{aligned} \quad (77)$$

If  $t_{com} = 2t_{calc}$  and  $J = M = N$ , then (77) can be simplified as follows:

$$T_{N_{proc}}^{(First)} = [(N^2 \times L) + (N^2 \times L + 3N_{proc} \times \log_2 N_{proc})] \times t_{calc} \quad , \quad (78)$$

then:

$$\begin{aligned} S_{N_{proc}}^{First} &= \frac{L^2 \times N^4}{L^2 \times N^4 + L \times N^2 + 3N_{proc} \times \log_2 N_{proc}} = \\ &= \frac{L \times N^2}{L \times N^2 + 3N_{proc} \times \log_2 N_{proc}} \quad . \end{aligned} \quad (79)$$

From (79) it follows that, if  $N$  and  $L$  are fixed and  $N_{proc}$  increases, speed-up degrades.

Substituting (66), (67) and (68) in (76) gives:

$$\begin{aligned}
T_{N_{proc}}^{(Second)} &= [N_{proc} \times ((L_{loc} + |E_p^q|) \times L'_{loc}) \times (J \times M) \times (N \times N) + (80) \\
&+ N_{proc} \times |E_p^q| \times (J \times M) \times |E_p^q| \times t_{calc} + \\
&+ [N_{proc} \times |E_p^q| \times (J \times M) \times |E_p^q| \times t_{com} \quad .
\end{aligned}$$

If  $t_{com} = 2t_{calc}$ , and  $J = M = N$ , and  $L = L'$ , since  $L_{loc}$  is at the most  $|E_p^q| = |J^q|$ , (80) can be simplified as follows:

$$T_{N_{proc}}^{(Second)} = [2 \times \frac{L^2}{N_{proc}} \times N^4 + 3 \frac{L^2}{N_{proc}} \times N^2] \times t_{calc} \quad , \quad (81)$$

then:

$$\begin{aligned}
S_{N_{proc}}^{Second} &= \frac{L^2 \times N^4}{2 \times \frac{L^2}{N_{proc}} \times N^4 + 3 \frac{L^2}{N_{proc}} \times N^2} = \\
&= \frac{N^2 \times N_{proc}}{2(N^2 + 3)} \quad . \quad (82)
\end{aligned}$$

Observe that if  $N$  and  $L$  are fixed and  $N_{proc}$  increases, also the speed-up increases, and the parallel algorithm scales with the processors number.

Second Strategy shows a better speed up. Other results are in [16].

## 6 The MEDITOMO library

In table 1 we list the software modules of the library <sup>(16)</sup>:

- *fully 3D*
  - *fan* collimator
    - \* *Conjugate Gradient* (`3dtomo_fan_cg`)
    - \* *Fixed Point* (`3dtomo_fan_fp+tv`)
    - \* *EMOSn* (`3dtomo_fan_em`)
    - \* *TV regularized EMOSn* (`3dtomo_fan_em+tv`)
  - *parallel* collimator

---

<sup>16</sup>`XXtomo_YYY_ZZ` is the name of each software:

- `XX` denotes the model (`3d` or `2d+1`);
- `YYY` denotes the geometry of the acquisition system (parallel `par` o fan `fan`);
- `ZZ` denotes the algorithm (conjugate gradient `cg`, expectation maximization `em` or fixed point `fp+tv`);

MEDITOMO Library	
Acquisition Models	
<i>Fully 3D</i>	<i>2D+1</i>
Reconstruction Algorithms	Reconstruction Algorithms
<ul style="list-style-type: none"> <li>* FP (TV)</li> <li>* CG</li> <li>* EMOS</li> <li>* EMOS-TV</li> </ul>	<ul style="list-style-type: none"> <li>* FP (TV)</li> <li>* CG</li> <li>* EMOS</li> </ul>

Table 1: MEDITOMO library

- \* *Conjugate Gradient* (3dtomo\_par\_cg)
- *2D+1*
  - *fan geometry*
    - \* *Conjugate Gradient* (2d+1tomo\_fan\_cg)
    - \* *Fixed Point* (2d+1tomo\_fan\_fp+tv)
    - \* *TV-based EMOSn* (2d+1tomo\_fan\_em+tv)

Parallel software use the **Message Passing Interface** (MPI) communication library for the message passing operations. We compare results obtained using:

1. 3dtomo\_fan\_cg and 3dtomo\_fan\_fp+tv;
2. 3dtomo\_fan\_em and 3dtomo\_fan\_em+tv;
3. 2d+1tomo\_fan\_em and 2d+1tomo\_fan\_cg.

Tests are performed on synthetic images and on real data:

- synthetic data obtained “*blurring*” the Hoffman phantom; in fig. 14, a projection for a fixed angle (a) and a sinogram (b) are shown. Projections data set is obtained simulating the acquisition system, where the matrix  $[K]_3$  is computed by using a Gauss function [9]. The values of (2) considered are  $\sigma_{min} = 3$ ,  $\sigma_{max} = 8$ ,  $r = 192$ . We generate 120 sinograms of size  $128 \times 128$ , where 120 is the number of angular views in  $[0, \pi)$ .
- in-vivo data<sup>(17)</sup>.

In figure 16 we show the *Restoration Error*:

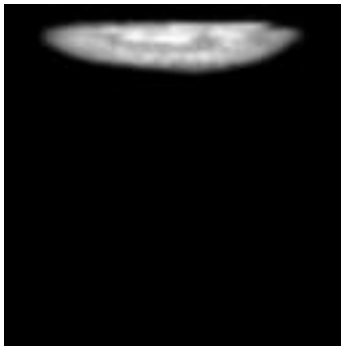
$$RE = \frac{\| [u]_{true} - [u^*] \|_2}{\| [u]_{true} \|_2} . \quad (83)$$

In figure 20 we show results for different iterations. Results refer to TV-based algorithms in figure 7 (a) and to the standard EMOSn (b).

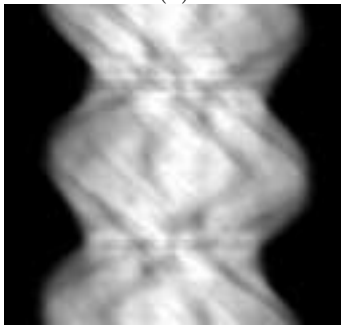
<sup>17</sup>The data are courteously granted by the Careggi Hospital in Florence (Italy).



Figure 13: *An Hoffman phantom section.*



(a)



(b)

Figure 14: A projection section (a) and a sinogram (b) examples

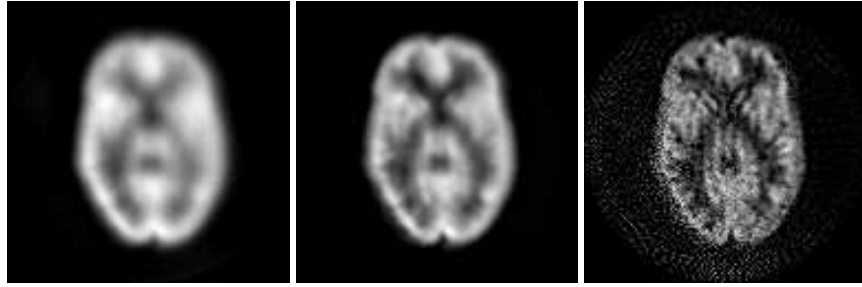


Figure 15: Results using `3dtomo_fan_cg` of the Hoffmann phantom. Number of steps are 7, 21, 63.

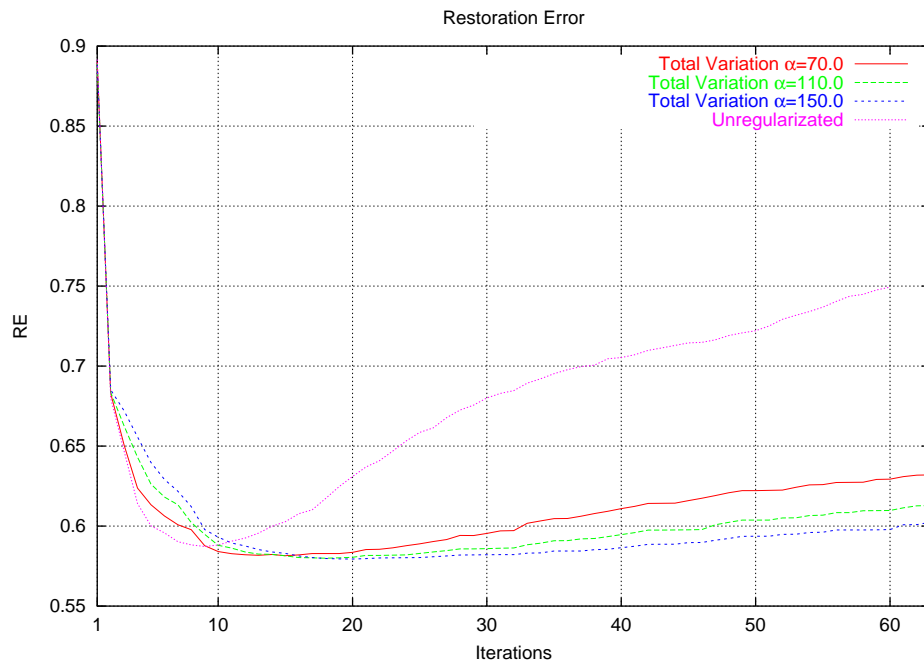


Figure 16:  $RE$  for `3dtomo_fan_cg` and `3dtomo_fan_fp+tv`. We plot  $RE$  at different values of the regularization parameter  $\alpha$ .

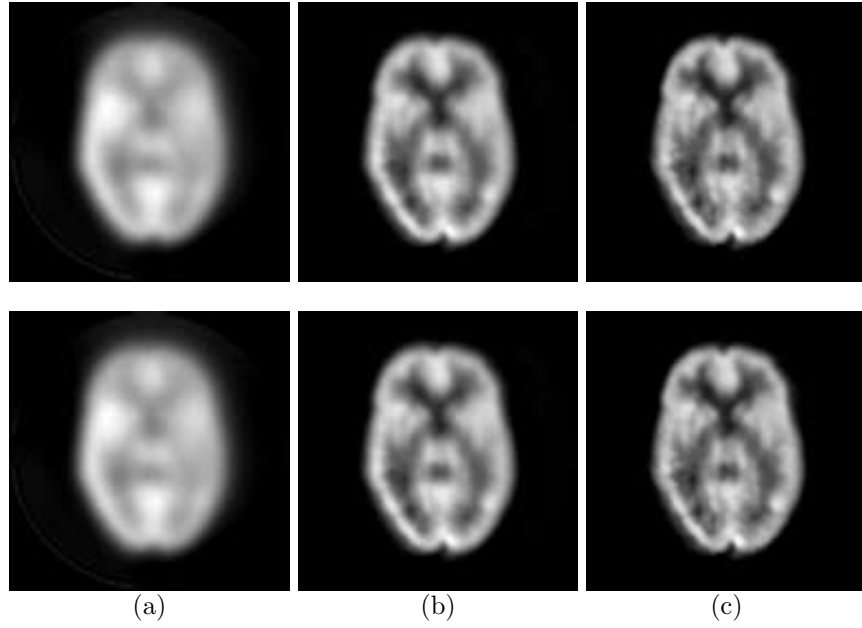


Figure 17: Results of `3dtomo_fan_fp+tv` on the Hofmann phantom, at step  $m = 1$  (a),  $m = 3$  (b) and  $m = 9$  (c) of FP. CG steps is 7 (the total number of iterations considered is than 7, 21 and 63 respectively).  $\alpha = 70$  at the first row and  $\alpha = 150$  at the second row.

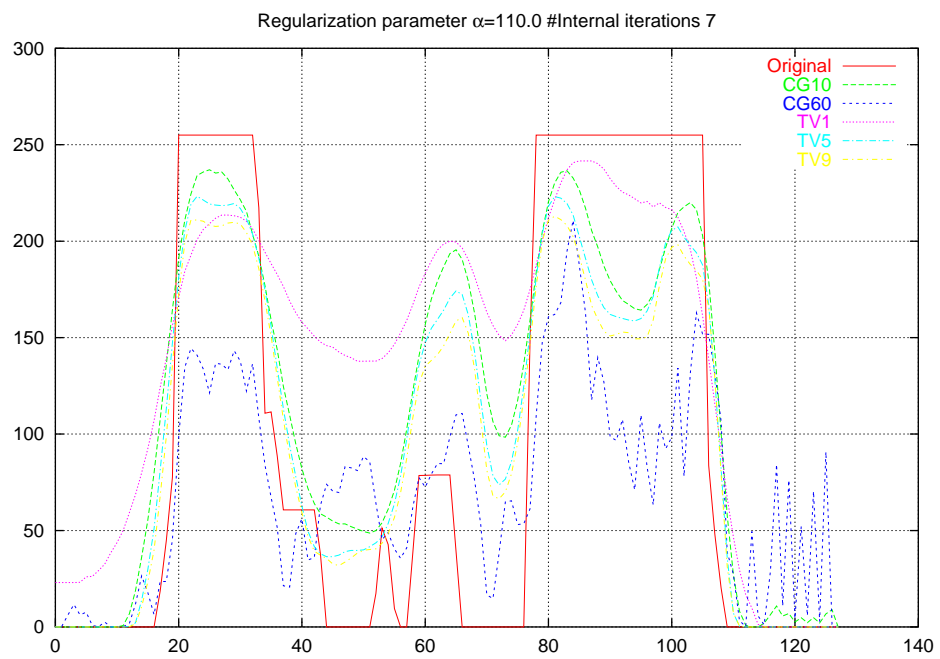


Figure 18: Profiles obtained by `3dtomo_fan_fp+tv` and `3dtomo_fan_cg`.



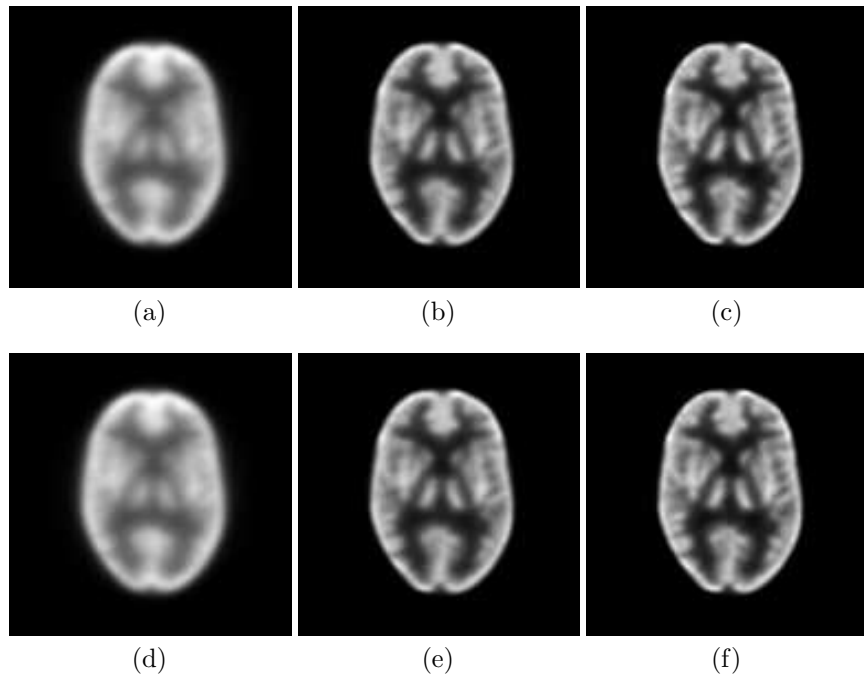
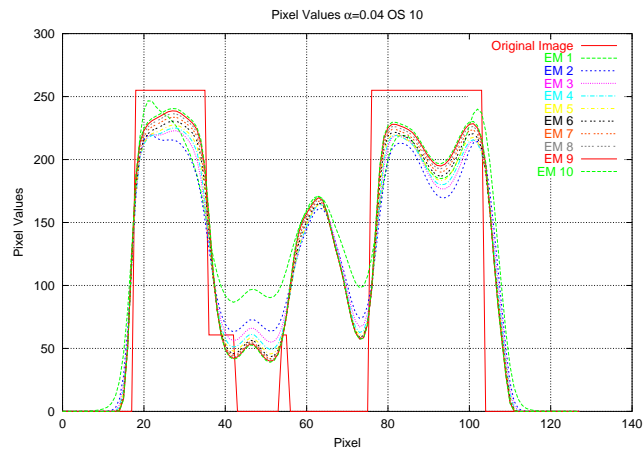
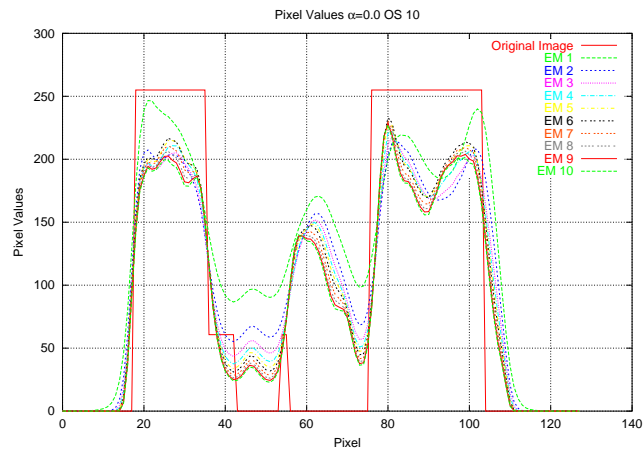


Figure 19: *Reconstruction of the Hoffman phantom, by 3dtomo\_fan\_em module, at I step (a), VI step (b) and X step (c) of EM. Reconstruction of Hoffman phantom, by 3dtomo\_fan\_em+tv, with  $\alpha = 0.04$ , at I step (d), VI step (e) and X step (f) of EM.*



(a)



(b)

Figure 20: Cross section of images obtained by `3dtomo_fan_em+tv` (a) and `3dtomo_fan_em` (b) respectively, for different steps of EM.

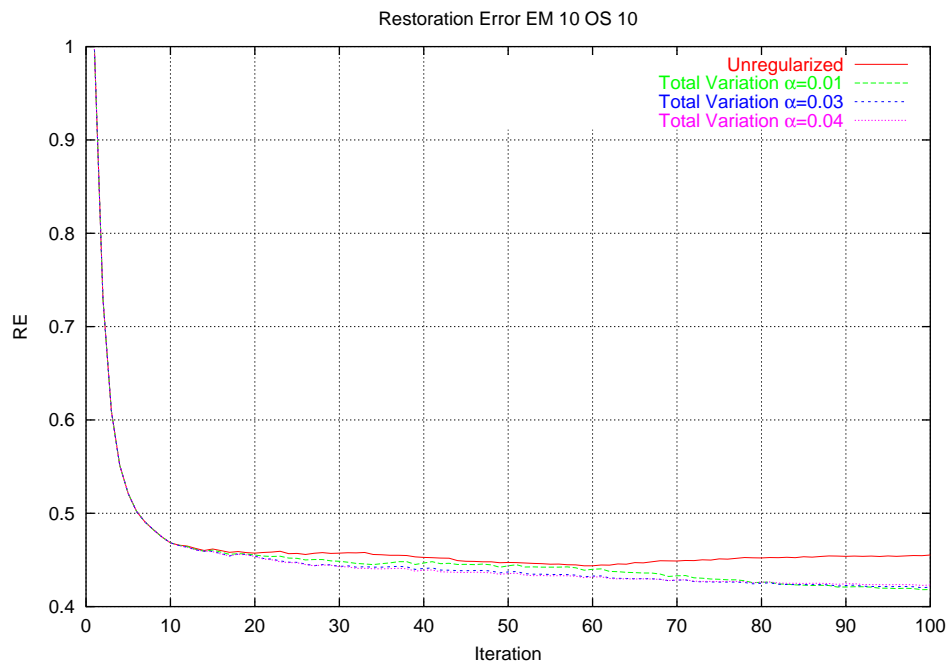


Figure 21:  $RE$  for  $2d+1$ tomo\_fan\_em and  $2d+1$ tomo\_fan\_em+tv.

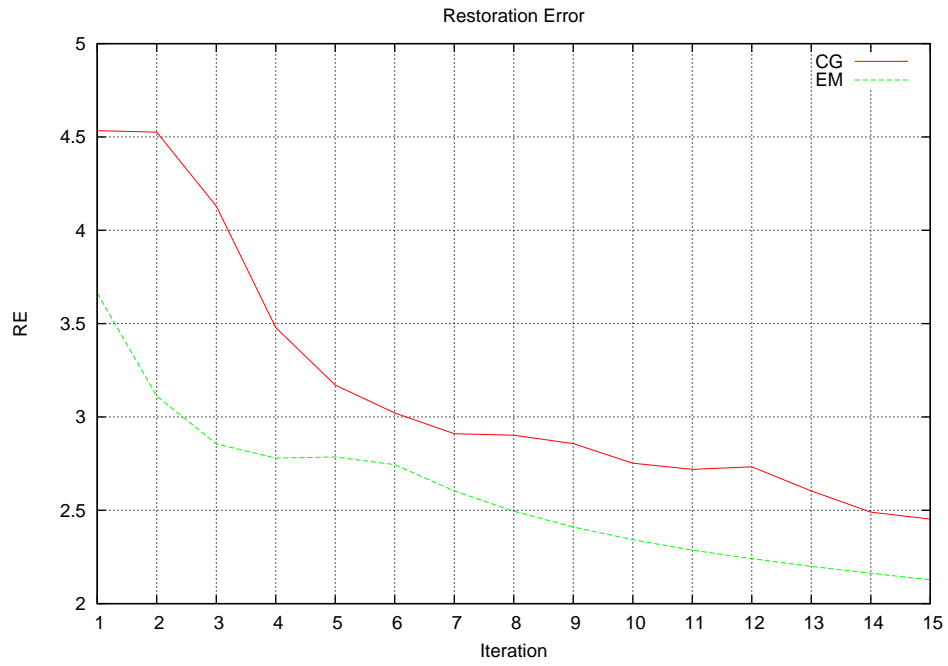


Figure 22: *RE* relatively to `2d+1tomo_fan_cg` and `2d+1tomo_fan_em` modules. The values of *RE* is plotted for different iterations of *OS*.

## 6.1 Conclusions

1. CG and EM appear to be *semiconvergent*: the RE, after reaching a minimum then grows while the TV-based algorithms show a more stable behaviour even with different values of  $\alpha$ .
2.  $it_{opt}$  is 10 and 6 for CG and EM respectively.
3. TV-regularization provides results of better quality.
4. Choice of  $\alpha$  and  $\beta$ :
  - 4.a Suitable values of  $\alpha$  are  $\alpha = 70$  for FP and  $\alpha = 0.04$  for EM;
  - 4.b parameter  $\beta$  in (??) has been fixed to 0.01 and 0.001 for the CG and EM. The value of  $\beta$  in EM is smaller because EM works with normalized values of the brightness.
5. EM is more accurate than CG. Accuracy increases as the number  $n$  of the *subset* grows. Optimal value of  $n$  is about 15-20 for EM and 10 for CG.
6. MEDITOMO has been tested on MIMD distributed memory parallel systems. In figures 24, 25 and 26 we show performance, in terms of time, speed-up and efficiency, of the `3dtomo_fan_fp+tv` running the software on a *Beowulf* system made of 18 nodes, each with 1.5 Ghz *Pentium IV* processor with a *Fast Ethernet 10/100 MBit* network card. Execution time of one iteration decreases from 10 to 1 minute. Efficiency of the most expensive code is at least of 60%.

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*The authors*

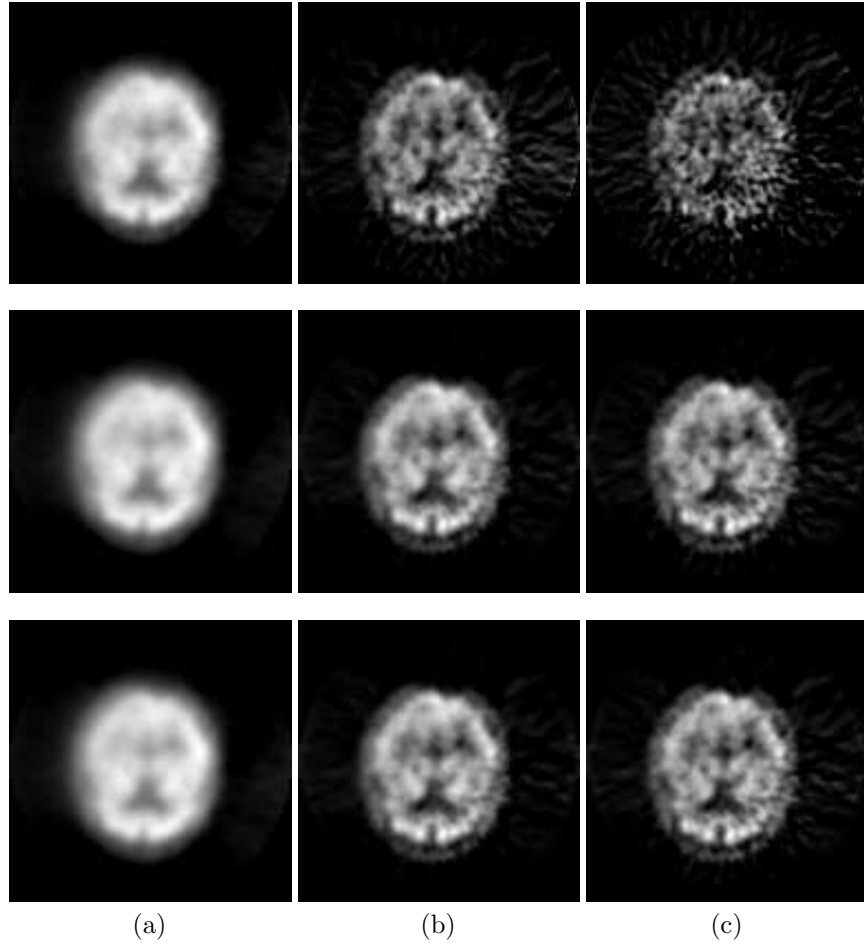


Figure 23: *First row: sections of human brain obtained by 3dtomo\_fan\_cg. iteration number: 5, 15 45 ; second and third rows: sections of human brain obtained by 3dtomo\_fan\_fp+tv with  $\alpha = 50,70$  at step  $m = 1$  (a),  $m = 3$  (b) and  $m = 5$  (c) of FP, and with CG steps 5 (Total number of iterations equals to 5, 15 45 respectively).*

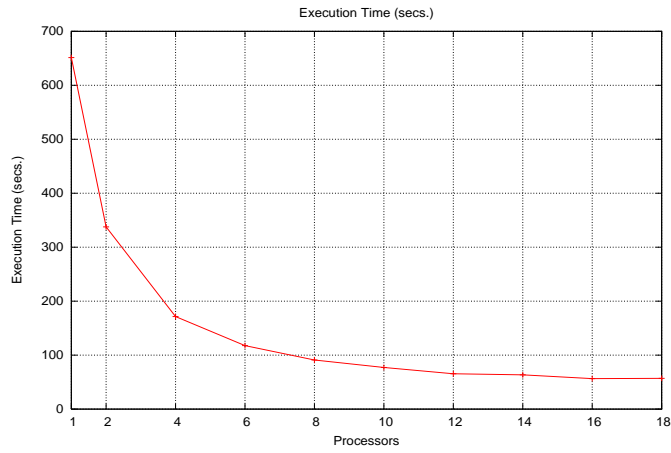


Figure 24: Execution time of 3dtomo\_fan\_fp+tv versus the number of processors.

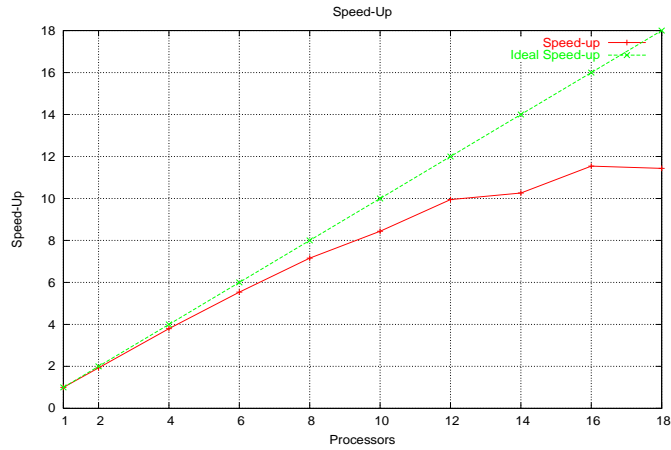


Figure 25: Speedup of 3dtomo\_fan\_fp+tv versus the number of processors.

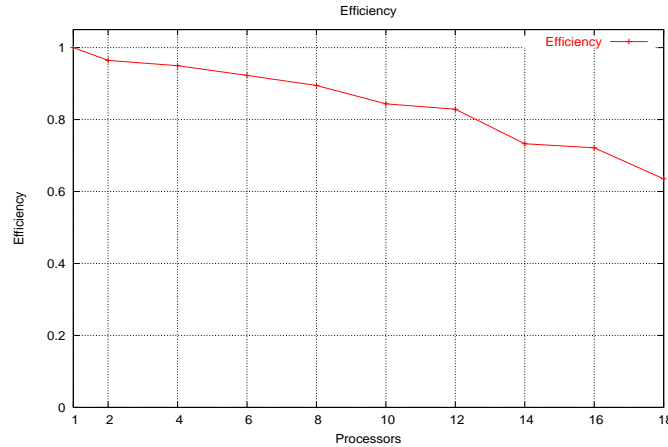


Figure 26: *Efficiency of 3dtomo\_fan\_fp+tv versus the number of processors.*

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